

Consensus in Systems with Eventual Quiescence

Brief Announcement

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SSS2025, October 11th, 2025

Motivation

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Paper: “*Reaching Consensus in the Presence of Contention-Related Crash Failures*”
by Anaïs Durand, Michel Raynal & Gadi Taubenfeld, SSS 2022

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Result:

Algorithm solving consensus tolerating up to k λ -constrained crashes, when $\lambda = n - k$, using objects with consensus number k

Consensus number of an object O is the maximal number of processes for which consensus can be solved despite any number of process crashes (occurring at any time) with any number of objects O and read/write registers

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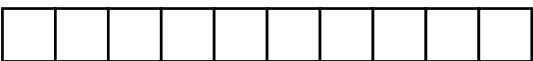
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Thank you!