

#### Stand Up Indulgent Rendezvous

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Slides available at https://bramas.fr





Two Robots



- Two Robots
- No Communication



- Two Robots
- No Communication
- No Memory



- Two Robots
- No Communication
- No Memory
- Disoriented



- Two Robots
- No Communication
- No Memory
- Disoriented
- No Common Unit-Distance



**ICU3E** 

3

What they do?
 Look (where is the other robot?)
 Compute (where am I going next?)

Move (who am I?)



3

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   Look (where is the other robot?)
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- When?



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- When?
  - Fully-Synchronous: Robot 1: LCM LCM LCM LCM LCM LCM ... Robot 2: LCM LCM LCM LCM LCM LCM ...
  - Semi-Synchronous:
     Robot 1: LCM LCM LCM LCM ...
     Robot 2: LCM LCM LCM LCM ...





4

• What we want:



4

- What we want:
  - Starting from an arbitrary configuration



- What we want:
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  - After a finite number of rounds



- What we want:
  - Starting from an arbitrary configuration
  - After a finite number of rounds
  - End up (and remain) at the same position



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5

• Suzuki-Yamashita (99)



5

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  - Solvable in FSYNC



5

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  - Solvable in FSYNC
  - Solvable in SSYNC with common coordinate system.



- Suzuki-Yamashita (99)
  - Solvable in FSYNC
  - Solvable in SSYNC with common coordinate system.
  - Not solvable in SSYNC without common coordinate system.





6

Two Possibilities:



6

- Two Possibilities:
  - Ignore crashed robots (weak gathering)



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- Two Possibilities:
  - Ignore crashed robots (weak gathering)
  - Gather at the position of the crashed robot (strong gathering)



**ICU3E** 

7

**ICU3E** 

7

• Strong gathering:



7

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  - Not possible when f > 1



7

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- Weak gathering (n-robots) tolerating crash faults:
  - possible for any *f* in SSYNC with multiplicity detection



#### The Stand Up Indulgent Rendezvous Problem

- What we want:
  - Starting from an arbitrary configuration
  - After a finite number of rounds
  - End up (and remain) at the same position
  - Even if one robot crashes (strong gathering)



#### Summary

	Rendezvous	SUIR
SSYNC oblivious, disoriented	Impossible [SY99]	
SSYNC oblivious, common x-axis		
SSYNC oblivious, common x-y-axis	Possible [SY99]	
SSYNC luminous, common x-y-axis	Possible	
FSYNC oblivious, disoriented	Possible [SY99]	



9

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**Lemma**: in an execution, if only one robot *r* is activated, then there is a round when *r* is dictated to move to the other robot.

**Proof**: Indeed, if the other robot is crashed, we know that in finite number of rounds *r* moves to the other robot.



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The scheduler is fair otherwise there exists an execution where only r (or r') is activated, without moving to the other robot, contradicting the previous Lemma.



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After each round the robots are not gathered.





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A rule defines the target for a set of views



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Examples:



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Example execution:



















ICUBE 18





18









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- Are on an oriented line
- Have a common unit distance
- Have rigid movements





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ICUBE 22

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6

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Generalize our results with *n* robots tolerating at most one fault



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Thank you for your attention

